

### SERVICE AND TRANSFER SELECTION FOR FREIGHTS IN A SYNCHROMODAL NETWORK

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Motivation

- Problem: service and transfer selection in a synchromodal network
- Proposed solution:
  - Markov Decision Process model
  - > Approximate Dynamic Programming algorithm
- Numerical experiments
- ••• What to remember



### **MOTIVATION** TRANSPORTATION OF CONTAINERS IN THE HINTERLAND



3/22



### MOTIVATION SYNCHROMODALITY





\*Source of video: Dutch Institute for Advanced Logistics (DINALOG) www.dinalog.nl UNIVERSITY OF TWENTE. 

# MOTIVATION

SYNCHROMODALITY



- The characteristics of synchromodality:
- Mode-free booking for all freights.
- Network-wise decisions at any point in time.
- Real-time information about the state of the network.
- Overall performance in the network and in time.





\*Source of artwork: European Container Terminals (ECT) – The future of freight transport (2011). UNIVERSITY OF TWENTE.

# SERVICE AND TRANSFER SELECTION IN A SYNCHROMODAL NETWORK

PROBLEM DESCRIPTION



# SERVICE AND TRANSFER SELECTION IN A SYNCHROMODAL NETWORK

A SOLUTION EXAMPLE

![](_page_6_Figure_2.jpeg)

![](_page_7_Picture_0.jpeg)

STOCHASTIC PROCESS UNDER SEQUENTIAL DECISION MAKING

- Stages for sequential decisions:  $t \in \mathcal{T} = \{0, 1, \dots, T^{\max} 1\}$
- Stochasticity in the arrival of freights:

Parameter	Notation	Probability
Number of freights	$f \in \mathcal{F} \subseteq \mathbb{N}$	$p_{f,t}^{\mathrm{F}}$
Origin	$i \in \mathcal{N}_t^{\mathcal{O}}$	$p_{i,t}^{\mathbf{O}}$
Destination	$d \in \mathcal{N}_t^{\mathrm{D}}$	$p_{d,t}^{\mathrm{D}}$
Release-day	$r \in \mathcal{R}_t = \{0, 1, 2, \dots, R_t^{\max}\}$	$p_{r,t}^{\mathrm{R}}$
Time-window length	$k \in \mathcal{K}_t = \{0, 1, 2, \dots, K_t^{\max}\}$	$p_{k,t}^{\mathrm{K}}$

- Decisions in which service to use for a freight, if any, at each stage.
- **Objective** to minimize a cost function over all stages.

![](_page_8_Picture_0.jpeg)

STATE AND DECISION AT EACH STAGE

The state describes all freights known at a stage:

$$S_t = [F_{i,d,r,k,t}]_{\forall i \in \mathcal{N}_t^{\mathcal{O}} \cup \mathcal{N}_t^{\mathcal{I}}, d \in \mathcal{N}_t^{\mathcal{D}}, r \in \mathcal{R}'_t, k \in \mathcal{K}_t}$$
(1)

The *decision* describes all freights assigned to the services at a stage:

$$x_t = [x_{i,j,d,k,t}]_{\forall (i,j) \in \mathcal{A}_t, d \in \mathcal{N}_t^{\mathrm{D}}, k \in \mathcal{K}_t}$$
(2a)

s.t.

$$\sum_{i \in \mathcal{N}_{t}^{\mathrm{I}} \cup \{d\}} x_{i,j,d,k,t} \leq F_{i,d,0,k,t}, \quad \forall i \in \mathcal{N}_{t}^{\mathrm{O}} \cup \mathcal{N}_{t}^{\mathrm{I}}, d \in \mathcal{N}_{t}^{\mathrm{D}}, k \in \mathcal{K}_{t}$$
(2b)

$$x_{i,j,d,k,t} = 0, \quad \forall (i,j) \in \mathcal{A}_t, d \in \mathcal{N}_t^{\mathrm{D}}, k \in \mathcal{K}_t | k < M_{i,j,t} + M_{j,d,t}$$
(2d)

$$\sum_{d \in \mathcal{N}_t^{\mathrm{D}}} \sum_{k \in \mathcal{K}_t} x_{i,j,d,k,t} \le Q_{i,j,t}, \quad \forall (i,j) \in \mathcal{A}_t^{\mathrm{I}}$$
(2e)

![](_page_9_Picture_0.jpeg)

EXOGENOUS INFORMATION AND TRANSITION FUNCTION

The exogenous information describes all freights that arrived between the previous and the current stage:

$$W_t = \left[\widetilde{F}_{i,d,r,k,t}\right]_{\forall i \in \mathcal{N}_t^{\mathcal{O}}, d \in \mathcal{N}_t^{\mathcal{D}}, r \in \mathcal{R}_t, k \in \mathcal{K}_t}$$
(3)

• The *transition function* describes how the system evolves::

$$S_t = S^M \left( S_{t-1}, x_{t-1}, W_t \right)$$
 (4a)

$$F_{t,i,d,0,k} = F_{t-1,i,d,0,k+1} - \sum_{j \in \mathcal{A}_t} x_{t-1,i,j,d,k+1} + F_{t-1,i,d,1,k}$$

$$+ \sum_{j \in \mathcal{A}_t \mid M_{j,i,t}=1} x_{t-1,j,i,d,k+M_{j,i,t}},$$

$$\forall i \in \mathcal{N}_t^{\mathrm{I}}, d \in \mathcal{N}_t^{\mathrm{D}}, k+1 \in \mathcal{K}_t$$

$$(4c)$$

![](_page_10_Picture_0.jpeg)

TRANSITION FUNCTION – A SMALL EXAMPLE

- The *release-day r* is relative to the current day *t*.
- The *time-window length k* is relative to the release-day *r*.
- Consider F<sub>i,d,r,k,t</sub> freights with k=4 sent from terminal i to terminal j using a service that lasts 2 days:

	t=7	t=8	t=9	t=10	t=11
	Monday	Tuesday	Wednesday	Thursday	Friday
i	<b>F</b> <sub>i,d,0,4,7</sub>				
j		<b>F</b> <sub>j,d,1,2,8</sub>	<b>F</b> <sub>j,d,0,2,9</sub>		
d					<b>F<sub>d,d,0,0,11</sub></b>

![](_page_11_Picture_0.jpeg)

**OBJECTIVE AND SOLUTION ALGORITHM** 

• The *objective* is to minimize the expected costs in the horizon:

$$\min_{\pi \in \Pi} \mathbb{E} \left[ \sum_{t \in \mathcal{T}} C_t \left( x_t^{\pi} \right) = \sum_{t \in \mathcal{T}} \sum_{(i,j) \in \mathcal{A}_t} \left( C_{i,j,t} \cdot \sum_{d \in \mathcal{N}_t^{\mathbf{D}}} \sum_{k \in \mathcal{K}_t} x_{i,j,d,k,t}^{\pi} \right) \middle| S_0 \right]$$
(5)

 The solution can be obtained using Bellman's principle of optimality and backward induction:

$$V_{t}(S_{t}) = \min_{\substack{x_{t}^{\pi} \in \mathcal{X}_{t} \\ \mathbf{A} \\ \mathbf{A} \\ \mathbf{A} \\ \mathbf{C}_{t}(x_{t}^{\pi}) + \sum_{\substack{\omega \in \Omega_{t+1} \\ \mathbf{A} \\$$

ALGORITHMIC APPROACH FOR SOLVING LARGE MARKOV MODELS.<sup>1</sup>

Algorithm 1 Approximate Dynamic Programming Solution Algorithm **Require:**  $\mathcal{T}, \mathcal{F}, \mathcal{G}, \mathcal{D}, \mathcal{R}, \mathcal{K}, [C_{\mathcal{D}'}]_{\forall \mathcal{D}' \subset \mathcal{D}}, B_d, Q, S_0, N$ 1: Initialize  $\bar{V}_t^0, \forall t \in \mathcal{T}$ 2:  $n \leftarrow 1$ 3: while n < N do 4:  $S_0^n \leftarrow S_0$ 5: for t = 0 to  $T^{max} - 1$  do 6:  $\hat{v}_t^n \leftarrow \min_{\boldsymbol{x}_t^n} \left( C\left(\boldsymbol{S}_t^n, \boldsymbol{x}_t^n\right) + \bar{V}_t^{n-1}\left( S^{M, \boldsymbol{x}}\left(\boldsymbol{S}_t^n, \boldsymbol{x}_t^n\right) \right) \right)$ 7: if t > 0 then  $\bar{V}_{t-1}^{n}(\boldsymbol{S}_{t-1}^{n,x*}) \leftarrow U^{V}(\bar{V}_{t-1}^{n-1}(\boldsymbol{S}_{t-1}^{n,x*}), \boldsymbol{S}_{t-1}^{n,x*}, \hat{v}_{t}^{n})$ 8: end i 9:  $\boldsymbol{x}_{t}^{n*} \leftarrow \arg\min_{\boldsymbol{x}_{t}^{n}} \left( C\left(\boldsymbol{S}_{t}^{n}, \boldsymbol{x}_{t}^{n}\right) + \bar{V}_{t}^{n-1} \left( S^{M, x}\left(\boldsymbol{S}_{t}^{n}, \boldsymbol{x}_{t}^{n}\right) \right) \right)$ 10:11:  $\boldsymbol{S}_{t}^{n,x*} \leftarrow S^{M,x}\left(\boldsymbol{S}_{t}^{n}, \boldsymbol{x}_{t}^{n*}\right)$ 12:  $\boldsymbol{W}_{t}^{n} \leftarrow \operatorname{RandomFrom}\left(\Omega\right)$  $\boldsymbol{S}_{t+1}^{n} \leftarrow S^{M}\left(\boldsymbol{S}_{t}^{n}, \boldsymbol{x}_{t}^{n*}, \boldsymbol{W}_{t}^{n}
ight)$ 13:end for 14: 15: end while 16: return  $[\bar{V}_t^N]_{\forall t \in \mathcal{T}}$ 

1. For a comprehensive explanation see Powell (2010) Approximate Dynamic Programming.

A GRAPHICAL REPRESENTATION OF THE ALGORITHM

![](_page_13_Figure_2.jpeg)

DECISIONS BASED ON THE VALUE FUNCTION APPROXIMATION

The expected future costs are explained, to a certain extent, by the post-decision state:

$$V_{t}(S_{t}) = \min_{x_{t}^{\pi} \in \mathcal{X}_{t}} \left( C_{t}(x_{t}^{\pi}) + \sum_{\omega \in \Omega_{t+1}} p_{\omega}^{\Omega_{t+1}} \cdot V_{t+1} \left( S^{M}(S_{t}, x_{t}^{\pi}, \omega) \right) \right), \forall S_{t} \in \mathcal{S} \quad (6)$$
$$V_{t}^{n}(S_{t}^{n}) = \min_{x_{t}^{\pi} \in \mathcal{X}_{t}} \left( C_{t}(x_{t}^{\pi}) + \overline{V}_{t}^{n}(S_{t}^{x,n}) \right) \quad (7)$$

For the value function approximation we use the traditional basis functions approach (i.e., weighted features of a postdecision state):

$$\overline{V}_{t}^{x,n}\left(S_{t}^{x,n}\right) = \sum_{a \in \mathcal{A}} \theta_{a,t}^{n} \phi_{a}\left(S_{t}^{x,n}\right) \tag{8}$$

![](_page_15_Picture_0.jpeg)

UPDATING THE VALUE FUNCTION APPROXIMATION

- In each iteration, we observe a realization of the future costs throughout the time horizon.
- We update (i.e., improve) the value function approximation using a recursive least squares (LSQ) method for non-stationary data method: Observed

Feature

**Optimization** Matrix

Error

DECISIONS BASED ON A SECOND TYPE OF VALUE FUNCTION APPROXIMATION

 A second type of value function approximation using basis functions and sampling future costs of a fixed heuristic:

$$\overline{V}_{t}^{x,n}\left(S_{t}^{x,n}\right) = \alpha \sum_{a \in \mathcal{A}} \theta_{a,t}^{n} \phi_{a}\left(S_{t}^{x,n}\right) + (1-\alpha) \overline{C}_{t}^{n}\left(S_{t}^{x,n}\right)$$
(10)

The *policy* resulting from both value function approximations is the function of the post-decision features with the weights of the last iteration:

$$x_t^{\pi} = \arg\min\left(C_t\left(x_t^{\pi}\right) + \sum_{a \in \mathcal{A}(S_t^x)} \theta_{a,t}^N \phi_a\left(S_t^x\right)\right)$$
(11)

![](_page_17_Picture_0.jpeg)

- **Explore** the value of our two ADP designs.
- Very small network in a 15 day horizon.
- Three time-window profiles: (1) long, (2) intermediate and (3) short.

![](_page_17_Figure_5.jpeg)

![](_page_18_Picture_0.jpeg)

- Simulation study using 10 different initial states per instance.
- Four different planning policies:
  - Benchmark: balance between consolidation and postponement using "savings" of intermodal services in the shortest path from origin to destination.
    - ADP 1: weighted features using traditional learning of weights.
  - ADP 2: weighted features using sampling while learning at early iterations.
  - Sampling: using Monte Carlo simulation to estimate future costs of all feasible decisions.

![](_page_19_Picture_0.jpeg)

### **Table 1.** Results for Instance $I_1$

State	Total	Benchmark		ADP 1		ADP 2		Sampling	
	Freights	Solution T	ime (s)	Solution	Time (s)	Solution	Time $(s)$	Solution	Time $(s)$
1	Λ	19991	0 09	_13.6%	20 0/	_33.0%	101 21	_12 20%	688 20
2	7	14684	0.94	-12.8%	52.06	-32.7%	96.67	-39.9%	1687.18
3	5	13042	0.92	-13.1%	31.68	-27.5%	81.46	-41.5%	827.15
4	6	13863	0.94	-12.3%	32.99	-25.9%	81.42	-39.0%	832.67
<b>5</b>	6	13863	0.91	-12.0%	108.62	-30.0%	111.21	-42.3%	1356.80
6	6	13863	0.94	-10.4%	102.12	-31.3%	67.58	-42.9%	1317.73
7	5	13042	0.94	-12.6%	40.26	-23.4%	<mark>81.99</mark>	-41.5%	893.59
8	4	12221	0.92	-14.7%	37.44	-25.0%	78.41	-38.9%	547.66
9	2	10579	0.94	-14.9%	31.72	-29.9%	45.13	-42.4%	611.18
10	5	13042	1.01	-11.2%	30.81	-32.9%	42.92	-40.6%	727.28

![](_page_20_Picture_0.jpeg)

**Table 2.** Average results for Instance  $I_2$  and  $I_3$ 

Instance	Benchmark		ADP 1		ADP 2		Sampling	
	ce Solution Time (s) S		Solution Time (s) S		olution Time (s) \$		Solution Time (s)	
$I_2 \\ I_3$	$\frac{11078}{12874}$	$\begin{array}{c} 0.88\\ 1.01 \end{array}$	-5.2% 2.9%	$10.52 \\ 3.19$	$-9.8\% \\ 0.4\%$	$13.89 \\ 2.31$	-31.2% -3.3%	$217.19 \\ 36.95$

![](_page_20_Figure_4.jpeg)

![](_page_21_Picture_0.jpeg)

![](_page_21_Picture_1.jpeg)

WHAT TO REMEMBER

We developed an MDP model and ADP algorithm to select services for freights in a synchromodal network that consider stochastic freight arrival and performance over time.

The policy performance of the weighted features is improved using sampling during the learning phase of the ADP algorithm.

For realistic networks, further research in the *value function* approximation and decision making within the ADP algorithm is necessary for guaranteeing the best policies.

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![](_page_22_Picture_1.jpeg)

### THANKS FOR YOUR ATTENTION! ARTURO E. PÉREZ RIVERA

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